ILLC Project Course in Statistical Learning Theory

Mathias Winther Madsen mathias.winther@gmail.com

Institute for Logic, Language, and Computation University of Amsterdam

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Random Products

Problem

You enter the casino with a capital of $C_0 = 1$. In each round,

- 1. your capital doubles with 75% probability, and
- 2. decreases by a factor of 10 with 25% probability.

Is this good news? What will happen after t = 1000 games?

$$C_{1000} = F_1 \times F_2 \times F_3 \times \cdots \times F_{1000}$$

e.g.,

$$C_{1000} = 2 \times 2 \times \frac{1}{10} \times 2 \times \cdots \times \frac{1}{10} \times 2.$$

Moment-Generating Functions

Definition

The **moment-generating function** of a random variable *X* is

$$G(r) = E[e^{rX}]$$

whenever these expected values are well-defined.



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Example

Suppose $X \sim \text{Bernoulli}(p)$, i.e.,

x
 1
 0

$$Pr(X = x)$$
 p
 $1 - p$

What is G(r)?

Moment-Generating Functions

Example

Suppose $X \sim \text{Geometric}(1/2)$, i.e.,

x
 1
 2
 3
 4
 ...

$$Pr(X=x)$$
 $1/2$
 $1/4$
 $1/8$
 $1/16$
 ...

What is G(r)?

(cf. en.wikipedia.org/wiki/St._Petersburg_paradox)

Problem

Let

$$S_t = X_1 + X_2 + \dots + X_t$$

be a sum of i.i.d. variables.

Which of the following two series diverge the fastest?

 $S_1, S_2, S_3, S_4, \ldots$ from $q, 2q, 3q, 4q, \ldots$

or

 $e^{S_1}, e^{S_2}, e^{S_3}, e^{S_4}, \dots$ from $e^q, e^{2q}, e^{3q}, e^{4q}, \dots$



Theorem (the exponential Markov bound)

$$\Pr\left\{S_t > qt\right\} \leq \left(\frac{G(r)}{e^{rq}}\right)^t.$$

Proof.

The Markov bound plus the fact that

$$G_{S_t}(r) = G_X(r)^t.$$

Theorem (The Chernoff Bound)

For a sum of t Bernoulli variables with parameter p,

$$\Pr\left\{\frac{S_t}{t} > q\right\} \leq \left(\left(\frac{p}{q}\right)^q \left(\frac{1-p}{1-q}\right)^{1-q}\right)^t$$



Proof.

$$\Pr\left\{\frac{S_t}{t} > q\right\} \leq \left(\frac{pe^r + 1 - p}{e^{rq}}\right)^t.$$



Theorem (The reverse Chernoff Bound)

For a sum of t Bernoulli variables with parameter p,

$$\Pr\left\{\frac{S_t}{t} < q\right\} \leq \left(\left(\frac{p}{q}\right)^q \left(\frac{1-p}{1-q}\right)^{1-q}\right)^t$$







Theorem (the unit interval Hoeffding bound)

All of the previous theorems are also true for arbitrary bounded variables with $0 \le X \le 1$ and E[X] = p.

Theorem (the general Hoeffding bound)

For a sum of bounded variables with $a \le X \le b$ and E[X] = p,

$$\Pr\{S_t > qt\} \leq \left(\left(\frac{p-a}{q-a}\right)^{q-a} \left(\frac{b-p}{b-q}\right)^{b-q} \right)^{t/(b-a)}$$